

# Computing Resolutions: How and Why

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Math 255 Algebraic Curves

# Who am I?

- Algebraic statistics

PhD mathematics (May 2011), advisor: Bernd Sturmfels

“Algebraic methods for evaluating integrals in Bayesian statistics”

<http://math.berkeley.edu/~shaowei/>

- Singular learning theory

Sumio Watanabe, resolution of singularities

- Machine learning

Collaboration with Stanford artificial intelligence lab (Andrew Ng)

Unsupervised deep learning

**Why?**

# Integral Asymptotics

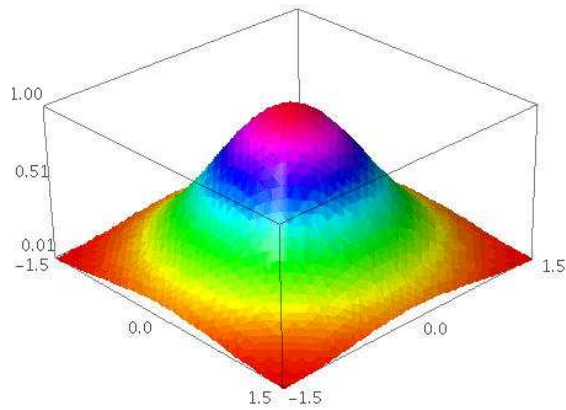
Stirling's approximation.

$$N! = N^{N+1} \int_0^\infty e^{-N(x-\log x)} dx \approx N^{N+1} \sqrt{\frac{2\pi}{N}} e^{-N}$$

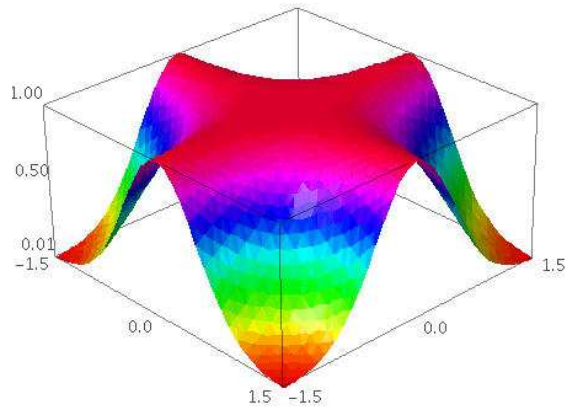
Laplace integrals.

$$Z(N) = \int_{\Omega} e^{-Nf(\omega)} \varphi(\omega) d\omega$$

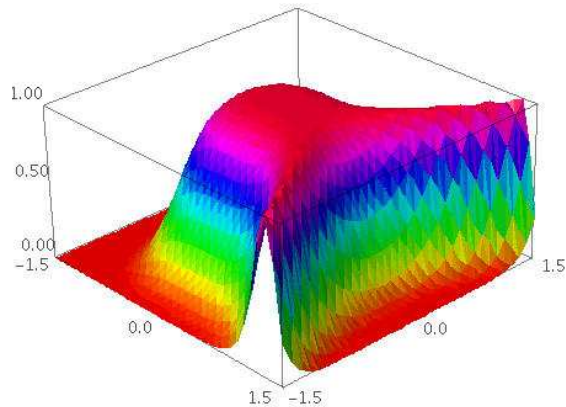
# Integral Asymptotics $Z(N) = \int_{\Omega} e^{-Nf(\omega)} \varphi(\omega) d\omega$



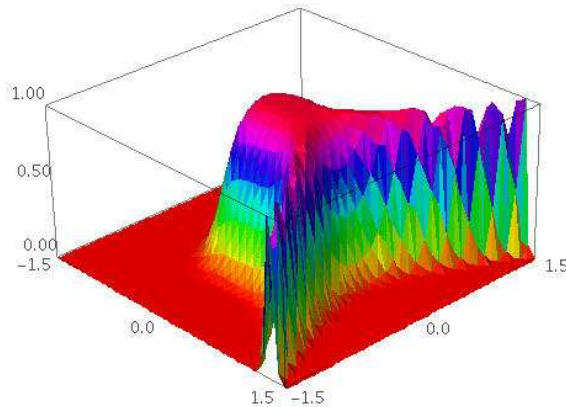
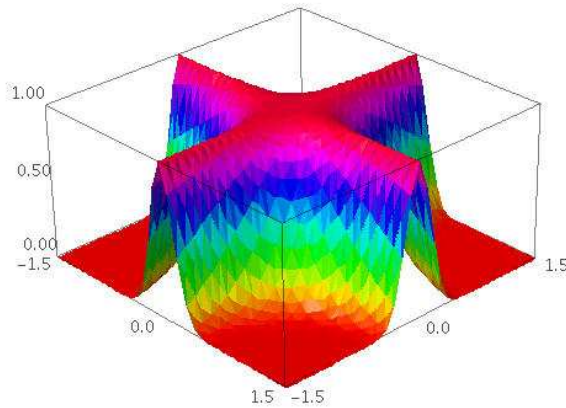
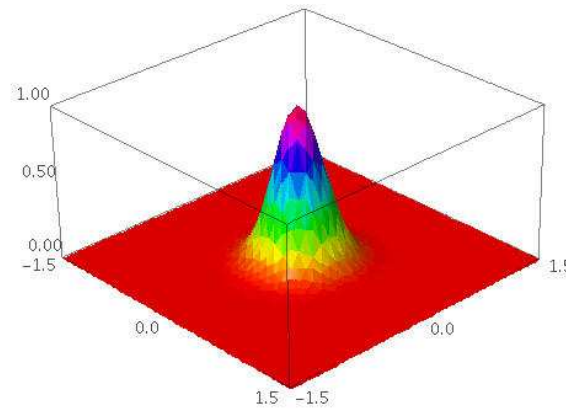
$$f(x, y) = x^2 + y^2$$



$$f(x, y) = (xy)^2$$



$$f(x, y) = (y^2 - x^3)^2$$



# Integral Asymptotics

Statistical integral.  $\int_{[0,1]^2} (1 - x^2 y^2)^{N/2} dx dy \approx$

$$\begin{aligned} & \sqrt{\frac{\pi}{8}} N^{-\frac{1}{2}} \log N & - \sqrt{\frac{\pi}{8}} \left( \frac{1}{\log 2} - 2 \log 2 - \gamma \right) N^{-\frac{1}{2}} \\ & - \frac{1}{4} N^{-1} \log N & + \frac{1}{4} \left( \frac{1}{\log 2} + 1 - \gamma \right) N^{-1} \\ & - \frac{\sqrt{2\pi}}{128} N^{-\frac{3}{2}} \log N & + \frac{\sqrt{2\pi}}{128} \left( \frac{1}{\log 2} - 2 \log 2 - \frac{10}{3} - \gamma \right) N^{-\frac{3}{2}} \\ & 0 & - \frac{1}{24} N^{-2} \quad + \dots \end{aligned}$$

Euler-Mascheroni constant  $\gamma = \lim_{n \rightarrow \infty} \left( \sum_{k=1}^n \frac{1}{k} - \log n \right) \approx 0.5772156649.$

# Log Canonical Thresholds

$I = \langle f_1, \dots, f_r \rangle \subset k[\omega_1, \dots, \omega_d]$  ideal

$k$  favorite field,  $W \subset k^d$  small nbd of origin.

*Complex LCT*  $\lambda$ : smallest pole of the zeta function

$$\zeta(z) = \int_W (|f_1(\omega)|^2 + \dots + |f_r(\omega)|^2)^{-z} d\omega, \quad z \in \mathbb{C}.$$

*Real LCT*  $\lambda$ : smallest pole of the zeta function

$$\zeta(z) = \int_W (f_1(\omega)^2 + \dots + f_r(\omega)^2)^{-z/2} d\omega, \quad z \in \mathbb{C}.$$

Independent of choice of generators!

*Real LCT*  $\lambda$ : coefficient in asymptotics

$$Z(N) = \int_W e^{-N(f_1(\omega)^2 + \dots + f_r(\omega)^2)} d\omega \approx C N^{-\lambda} (\log N)^{\theta-1}.$$

**How?**



# Blowups: Definition

Let  $Z \subset \mathbb{A}^d$  be a smooth variety whose ideal is  $\langle f_1, \dots, f_r \rangle$ .  
e.g.  $Z$  is the origin, whose ideal is  $\langle x_1, \dots, x_d \rangle$ .

Let us blow up  $\mathbb{A}^d$  with *center*  $Z$ .

For each point  $x = (x_1, \dots, x_d) \in \mathbb{A}^d$  not in  $Z$ , let us tag it with the point  $f^{\mathbb{P}}(x) = (f_1(x) : \dots : f_r(x)) \in \mathbb{P}^{r-1}$ . The set  $X$  of points

$$(x, f^{\mathbb{P}}(x)) \in \mathbb{A}^d \times \mathbb{P}^{r-1}, \quad x \in \mathbb{A}^d \setminus Z$$

has a Zariski closure  $\tilde{X}$  called the *blowup* of  $\mathbb{A}^d$  with center  $Z$ .

The projection  $\pi : \tilde{X} \subset \mathbb{A}^d \times \mathbb{P}^{r-1} \rightarrow \mathbb{A}^d$  is the *blowup map*.  
This map restricts to an isomorphism  $X \rightarrow \mathbb{A}^d \setminus Z$ , while the preimage  $E = \pi^{-1}Z \simeq Z \times \mathbb{P}^{\text{codim}(Z)-1}$  is the *exceptional divisor*.

# Blowups: Properties

- $\tilde{X}$  is smooth and is the union  $X \cup E$ .

- $\tilde{X}$  is the variety

$$\{(x, y) \in \mathbb{A}^d \times \mathbb{P}^{r-1} : y_i f_j(x) = y_j f_i(x) \text{ for all } i, j\}.$$

If the center  $Z$  is the origin, then  $\tilde{X}$  is a **toric** variety.

- $\tilde{X}$  is covered by affine charts  $U_i = \{(x, y) \in X : y_i \neq 0\}$  where

$$U_i \simeq \operatorname{Spec} k[x_1, \dots, x_d, \frac{f_1(x)}{f_i(x)}, \dots, \frac{f_r(x)}{f_i(x)}].$$

If the center  $Z$  is the origin, then  $U_i \simeq \mathbb{A}^d$  and it has coordinates

$$(\xi_1, \dots, \xi_d) = (\frac{y_1}{y_i}, \dots, x_i, \dots, \frac{y_d}{y_i}).$$

The blowup map  $\pi$  restricts to the affine map  $\pi_i : U_i \rightarrow \mathbb{A}^d$ ,

$$(\xi_1, \dots, \xi_d) \mapsto (\xi_1 \xi_i, \dots, \xi_i, \dots, \xi_d \xi_i).$$

# Blowups: Transforms

Let  $V \subset \mathbb{A}^d$  be a variety which contains a smooth subvariety  $Z$ .  
Let  $\pi : \tilde{X} \rightarrow \mathbb{A}^d$  be the blowing up of  $\mathbb{A}^d$  with center  $Z$ .

*Total* transform:  $\pi^{-1}(V)$

*Strict* (or proper) transform:  $\pi_{\text{st}}^{-1}(V) = \overline{\pi^{-1}(V \setminus Z)}$ .

Exceptional divisor:  $E = \pi^{-1}(Z)$

$$\pi^{-1}(V) = \pi_{\text{st}}^{-1}(V) \cup E$$

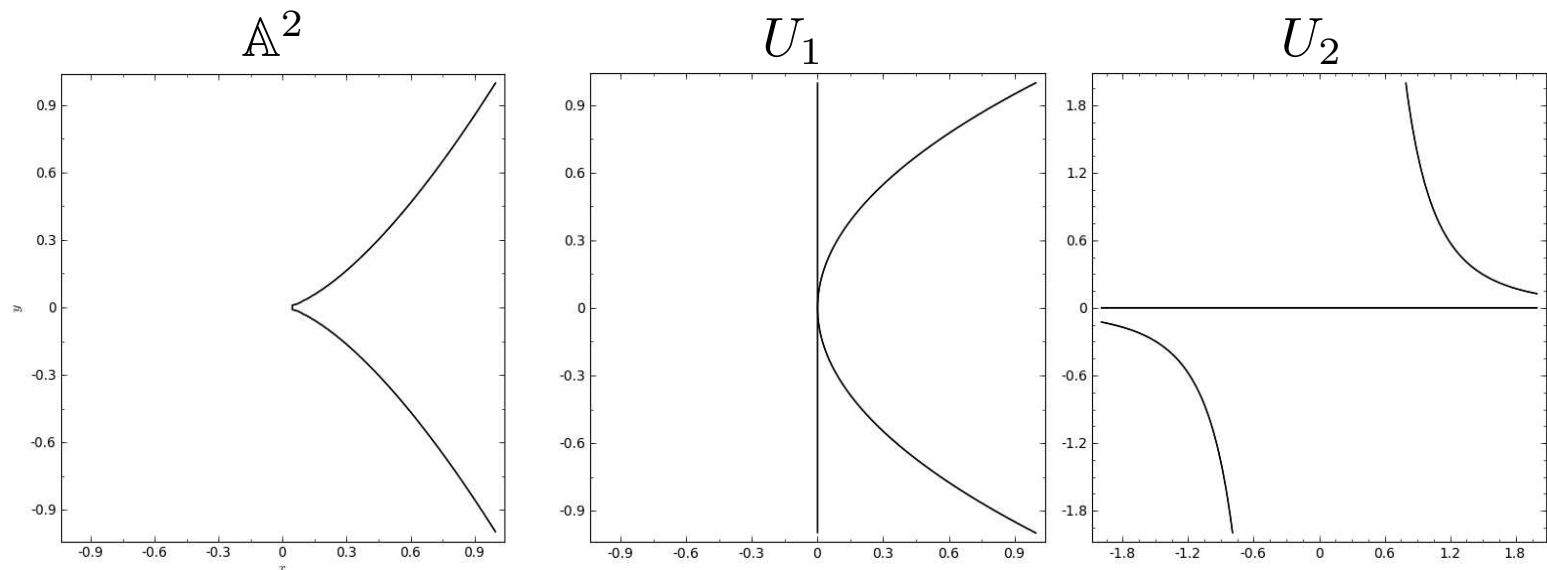
# Blowups: Example

Suppose we blow up the origin in  $\mathbb{A}^2$ . The blowup  $\tilde{X}$  can be thought of as a *mobius strip* with infinite width. It has two charts  $U_1, U_2$ .

Consider the cusp  $\mathcal{V}(y^2 - x^3) \subset \mathbb{A}^2$ . Under the blowing up,

●  $\pi_1 : U_1 \rightarrow \mathbb{A}^2, (s, t) \mapsto (x, y) = (s, st). \quad y^2 - x^3 = s^2(t^2 - s)$

●  $\pi_2 : U_2 \rightarrow \mathbb{A}^2, (u, v) \mapsto (x, y) = (uv, v). \quad y^2 - x^3 = v^2(1 - u^3v)$



**Software**

# Singular Code

```
LIB "sing.lib";
LIB "primdec.lib";
LIB "resolve.lib";
LIB "reszeta.lib";
LIB "resgraph.lib";

ring RR = 0, (x,y), dp;
ideal I = y^2-x^3;
ideal Z = x,y;
list L = blowUp(I,Z);
size(L);
def Q = L[1]; setring Q;
showBO(BO); // BO = Basic Object
setring RR;
```

# Singular Code

```
// resolve: "A" = All charts,  
//           "L" = Locally at origin
```

```
ideal I = y^2-x^3;  
list L = resolve(I, 0, "A", "L");  
presentTree(L);
```

```
ideal I = x^6+y^6-x*y;  
list L = resolve(I, 0, "A", "L");  
presentTree(L);
```

```
ideal I = x^6+y^6+y^2-x^3;  
list L = resolve(I, 0, "A", "L");  
presentTree(L);
```

# Resolution of Singularities

## For curves:

Blow up singularities (points) one by one.

## For varieties of higher dimension:

Blow up smooth subvarieties of singular locus.

Not easy to identify these smooth centers in general.

Knowing how to do this will get you a Field's Medal.

Dimension reduction via *hypersurfaces of maximal contact*



# Other Cool Facts

## ● Blowing up a linear subspace

Let  $\pi : \tilde{X} \rightarrow \mathbb{R}^n$  be the blowing up of the origin in  $\mathbb{R}^n$ . Then,  $\pi \times \text{id} : \tilde{X} \times \mathbb{R}^m \rightarrow \mathbb{R}^n \times \mathbb{R}^m$  is the blowing up of  $\{0\} \times \mathbb{R}^m$  in  $\mathbb{R}^{n+m}$ .

## ● Global blowups is the gluing of local blowups

Suppose we want to blow up a smooth center  $Z$  in  $\mathbb{R}^n$ . Cover  $\mathbb{R}^n$  with small affine charts, and pick coordinates in each chart so that  $Z$  is a linear subspace of each chart. Then, blowing up  $Z$  in  $\mathbb{R}^n$  is equivalent to blowing up the corresponding linear subspaces in each chart and gluing these maps together.

**Moral of the Story:** We only need to know how to blow up the origin!