

# Formal AI-Assisted Code Specification And Synthesis

## Concrete Steps Towards Safe Sociotechnical Systems



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Joint work with [Atlas Computing](#)



Atlas Computing

**Challenge**

# Safety and Security of Sociotechnical Systems

Intentional attacks/  
accidental mistakes

Human-led/  
AI-enabled

MICROSOFT / TECH / OPENAI

## Microsoft and OpenAI say hackers are using ChatGPT to improve cyberattacks



Photo by Amelia Holowaty Krales / The Verge

/ A number of nation-backed groups are starting to use large language models to help with research, scripting, and phishing emails.

By **Tom Warren**, a senior editor covering Microsoft, PC gaming, console, and tech. He founded WinRumors, a site dedicated to Microsoft news, before joining The Verge in 2012.

Feb 14, 2024, 4:00 AM PST



6

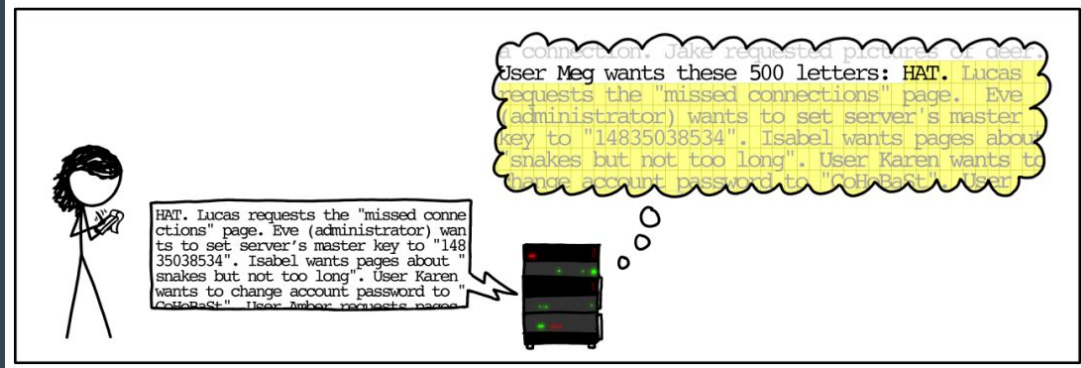
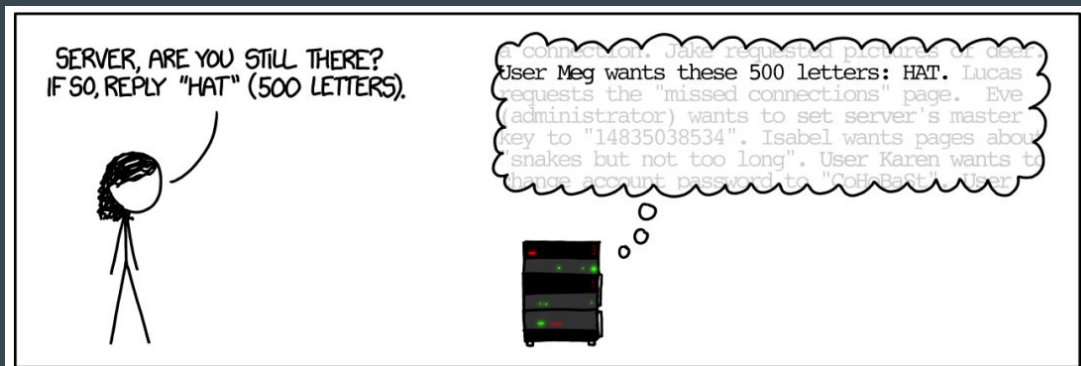
Comments (6 New)

If you buy something from a Verge link, Vox Media may earn a commission. [See our ethics statement.](#)

<https://www.theverge.com/2024/2/14/24072706/microsoft-openai-cyberattack-tools-ai-chatgpt>

# Hardening Systems with Formal Verification

Do we want to fix our sociotechnical systems one bug at a time?

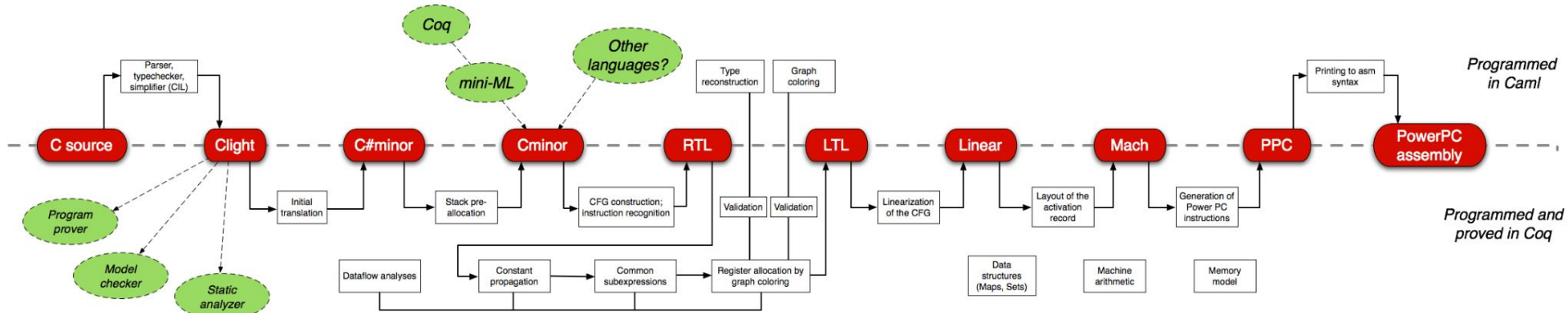


# Hardening Systems with Formal Verification

“.. **CompCert** is the only compiler we have tested for which Csmith cannot find wrong-code errors. This is not for lack of trying: we have devoted about six CPU-years to the task.”

Yang, X., Chen, Y., Eide, E. and Regehr, J., 2011, June. Finding and understanding bugs in C compilers. In Proceedings of the 32nd ACM SIGPLAN conference on Programming language design and implementation (pp. 283-294).

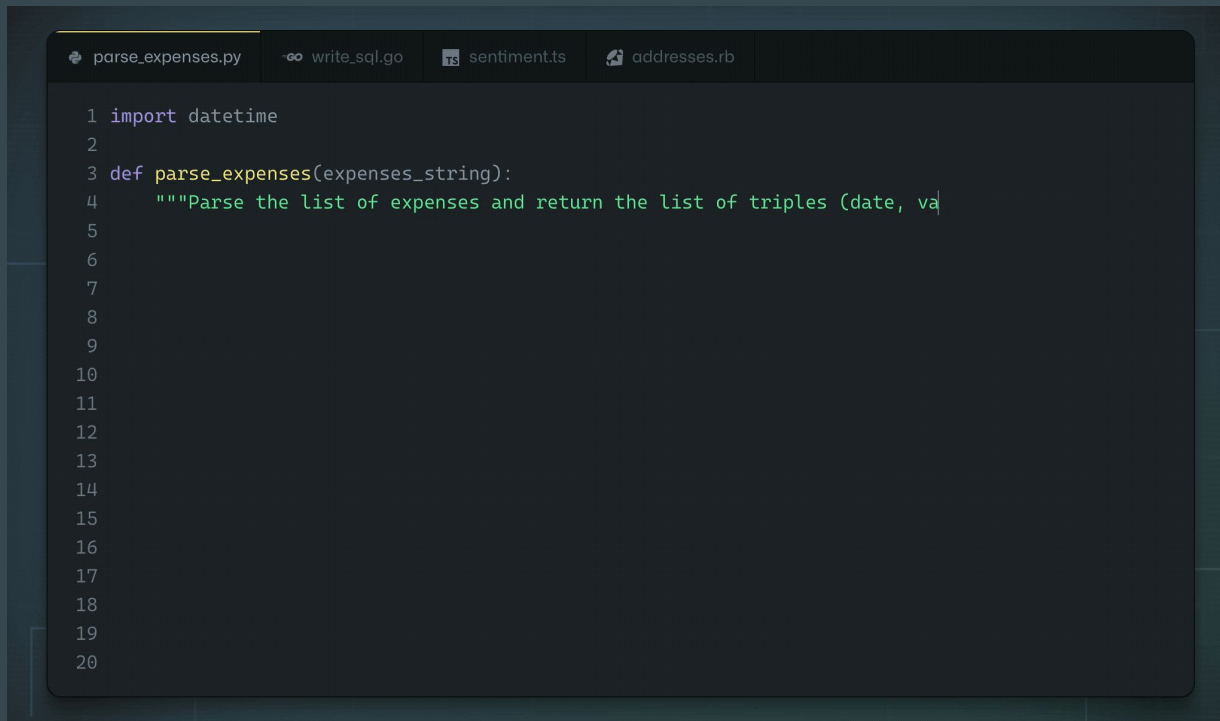
But writing formal specifications, implementations and proofs is hard!



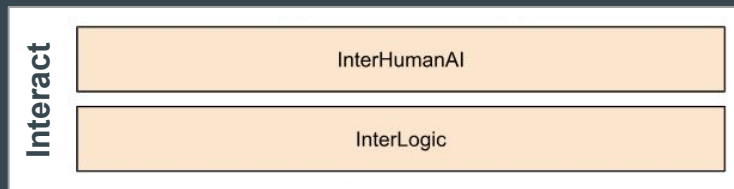
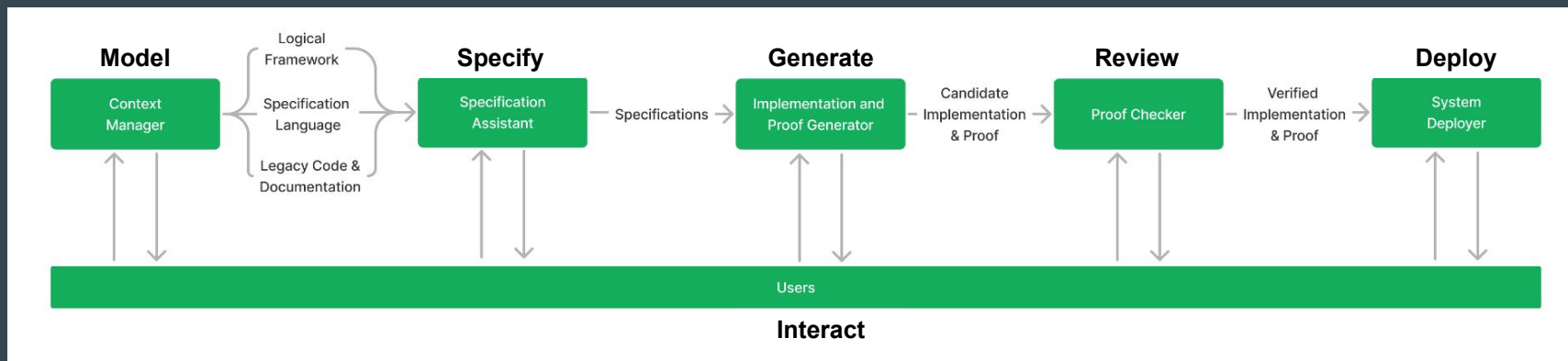
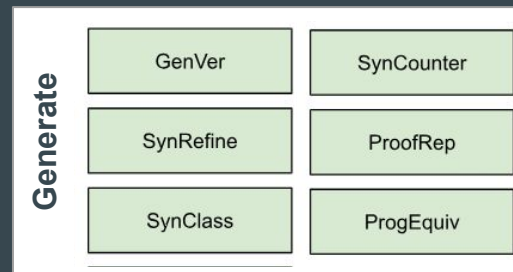
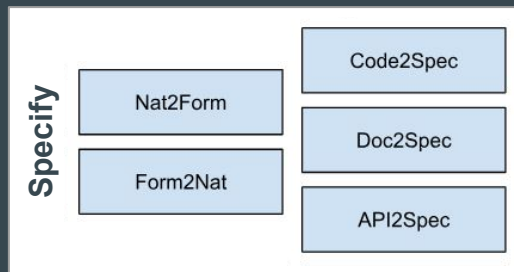
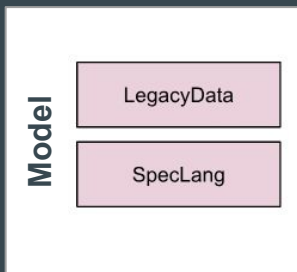
# Scaling FV with AI Assistance

Large language models (LLMs) are becoming good at generating code.

How do we ensure they produce code that is correct?

A screenshot of a code editor interface with a dark theme. The editor has four tabs at the top: 'parse\_expenses.py', 'write\_sql.go', 'sentiment.ts', and 'addresses.rb'. The active tab is 'parse\_expenses.py'. The code is written in Python and includes a docstring. The code is as follows:

```
1 import datetime
2
3 def parse_expenses(expenses_string):
4     """Parse the list of expenses and return the list of triples (date, va
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
```



# Roadmap

**Model**

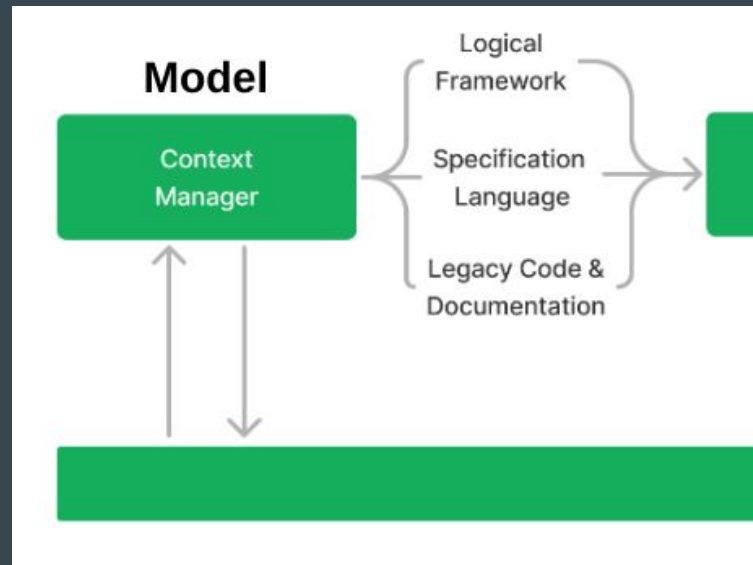


## Legacy Code and Docs (LegacyData)

- Verify legacy code, fix its bugs?
- Provide data/prompt to AI models for code specification and synthesis!

## Specification Languages (SpecLang)

- Domain specific languages



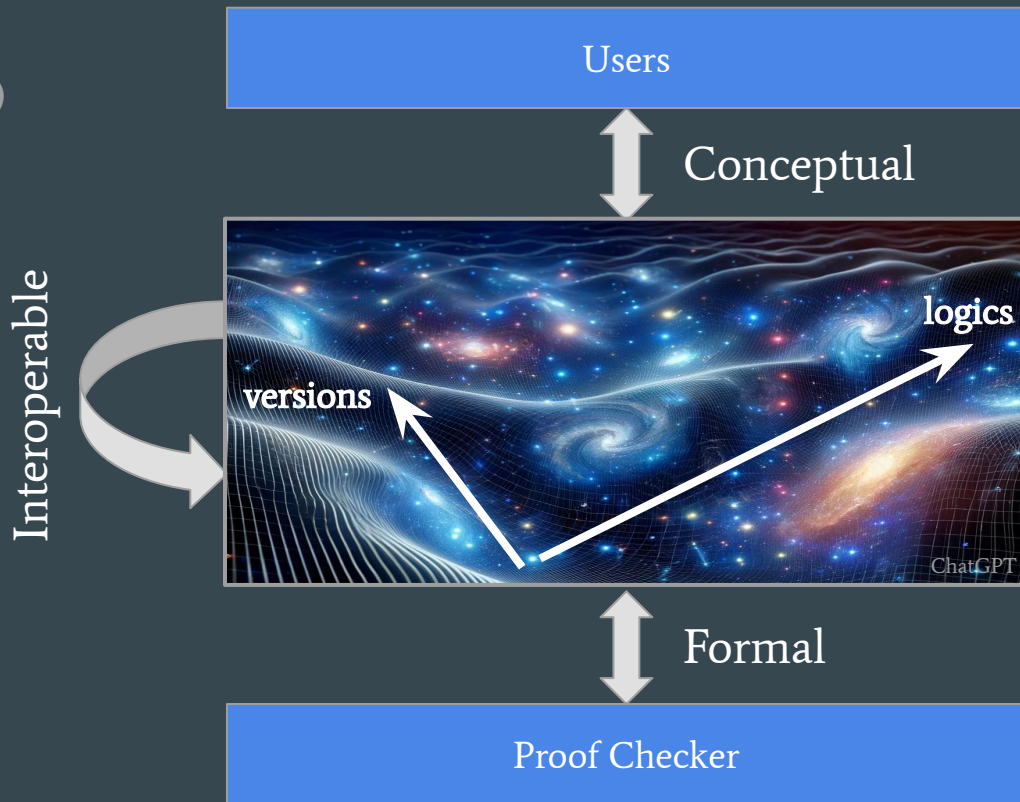
# Collective Modelling

(not in roadmap; work at Topos)

Double-categorical  
systems theory

Domain-specific  
logics (DSLs)

- Conceptual
- Interoperable
  - Between models
  - Between versions
- Formal



# Collective Sense-making

(prerequisite for roadmap)

Not top-down decision-making.

In a society saturated with top-down decision structures, building aligned AI models is key.

A better long-term solution is to rebuild society with a high value for collective sense-making, enabled by its sociotechnical systems.



**Interact**

## Human-AI Interactions

(InterHumanAI)

Human users

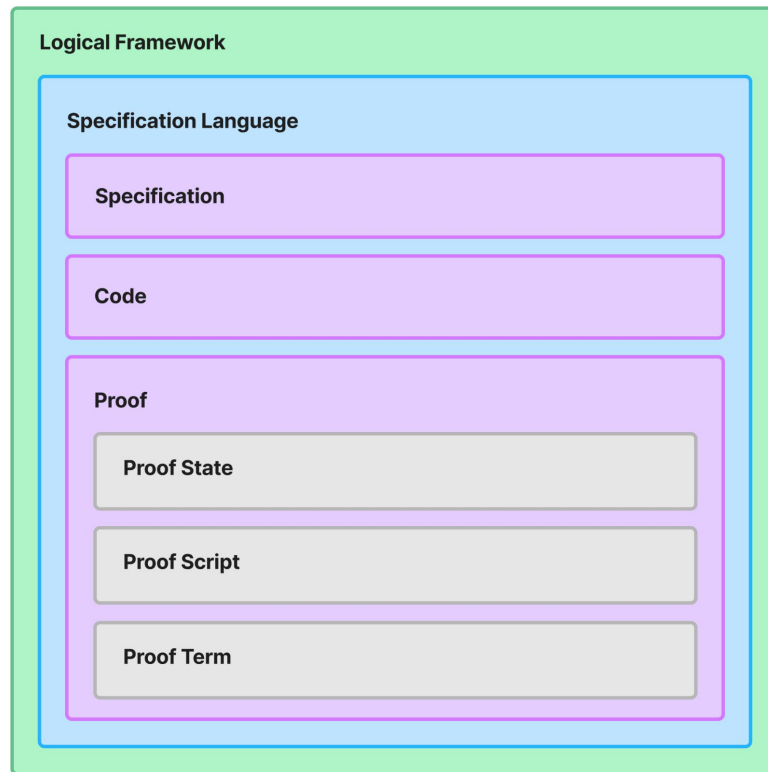
+ Proof assistant

+ AI assistant

Integration with Lean,

Frama-C, Verus, etc.

[Topos: ambitious program to  
redefine “proof assistant”]



# Logic Translations

## (InterLogic)



You

Translate the following code from the Coq programming language into the Lean programming language.

```

Definition instrDenote (i : instr) (s : stack) : option stack :=
  match i with
  | iConst n => Some (n :: s)
  | iBinop b =>
    match s with
    | arg1 :: arg2 :: s' => Some ((binopDenote b) arg1 arg2 :: s')
    | _ => None
  end
end
end.

```



ChatGPT

To translate the provided Coq code into the Lean programming language, we need to understand the structure of the Coq code and then replicate that logic in Lean. The Coq code defines a function `instrDenote` which takes an instruction `i` and a stack `s`, and then processes the instruction to potentially modify the stack.

Here's how the same functionality can be implemented in Lean using a similar pattern:

lean

Copy code

```

inductive instr
| iConst : ℕ → instr
| iBinop : (ℕ → ℕ → ℕ) → instr

open instr

def binopDenote (b : ℕ → ℕ → ℕ) (arg1 arg2 : ℕ) : ℕ :=
  b arg1 arg2

def instrDenote : instr → list ℕ → option (list ℕ)
| (iConst n) s := some (n :: s)
| (iBinop b) (arg1 :: arg2 :: s) := some (binopDenote b arg1 arg2 :: s)
| (iBinop _) _ := none

```

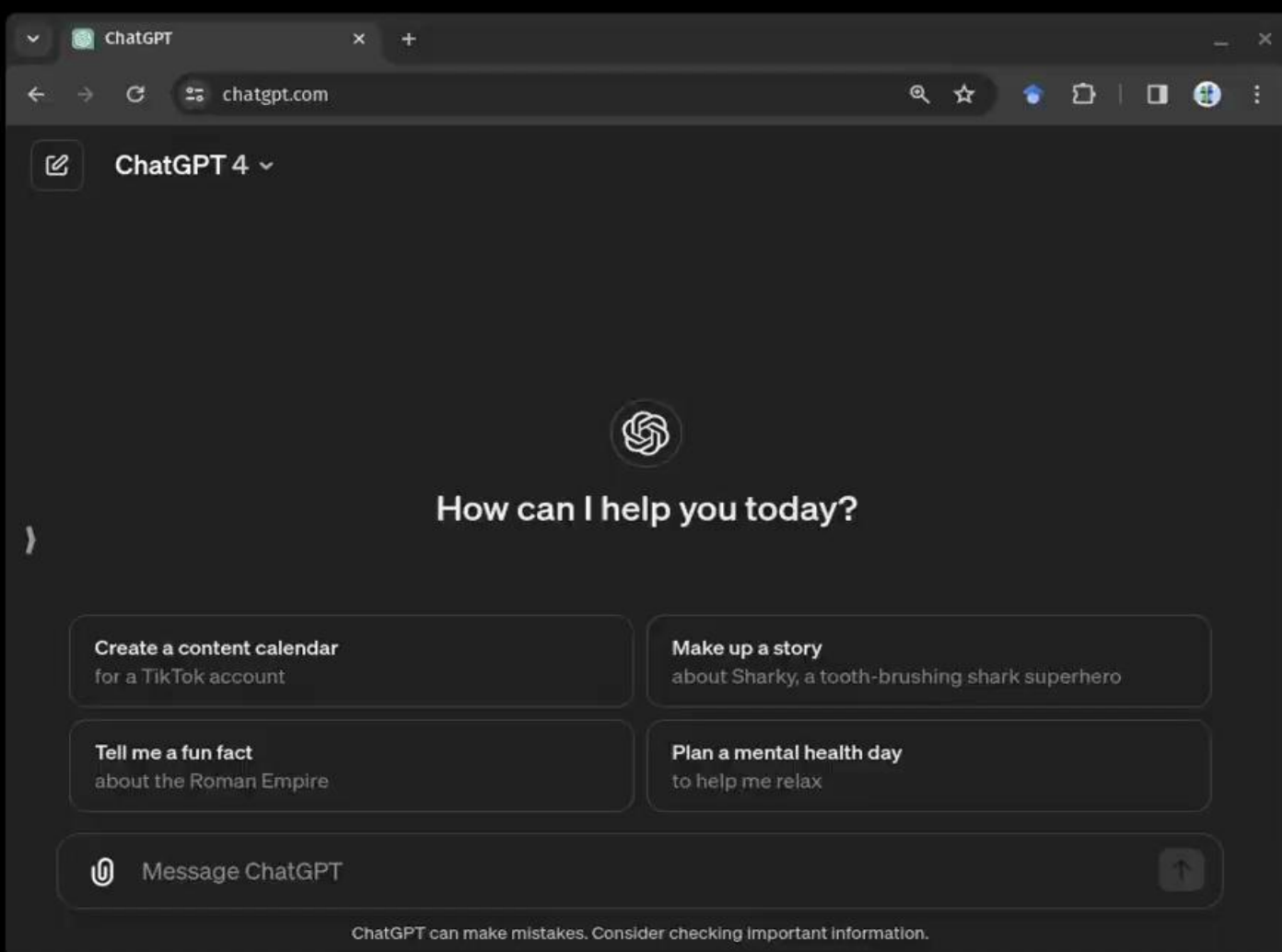
**Specify**

Specify

# From Natural Language To Specs

Importance of  
legacy docs

Generating  
well-formed  
terms of DSLs

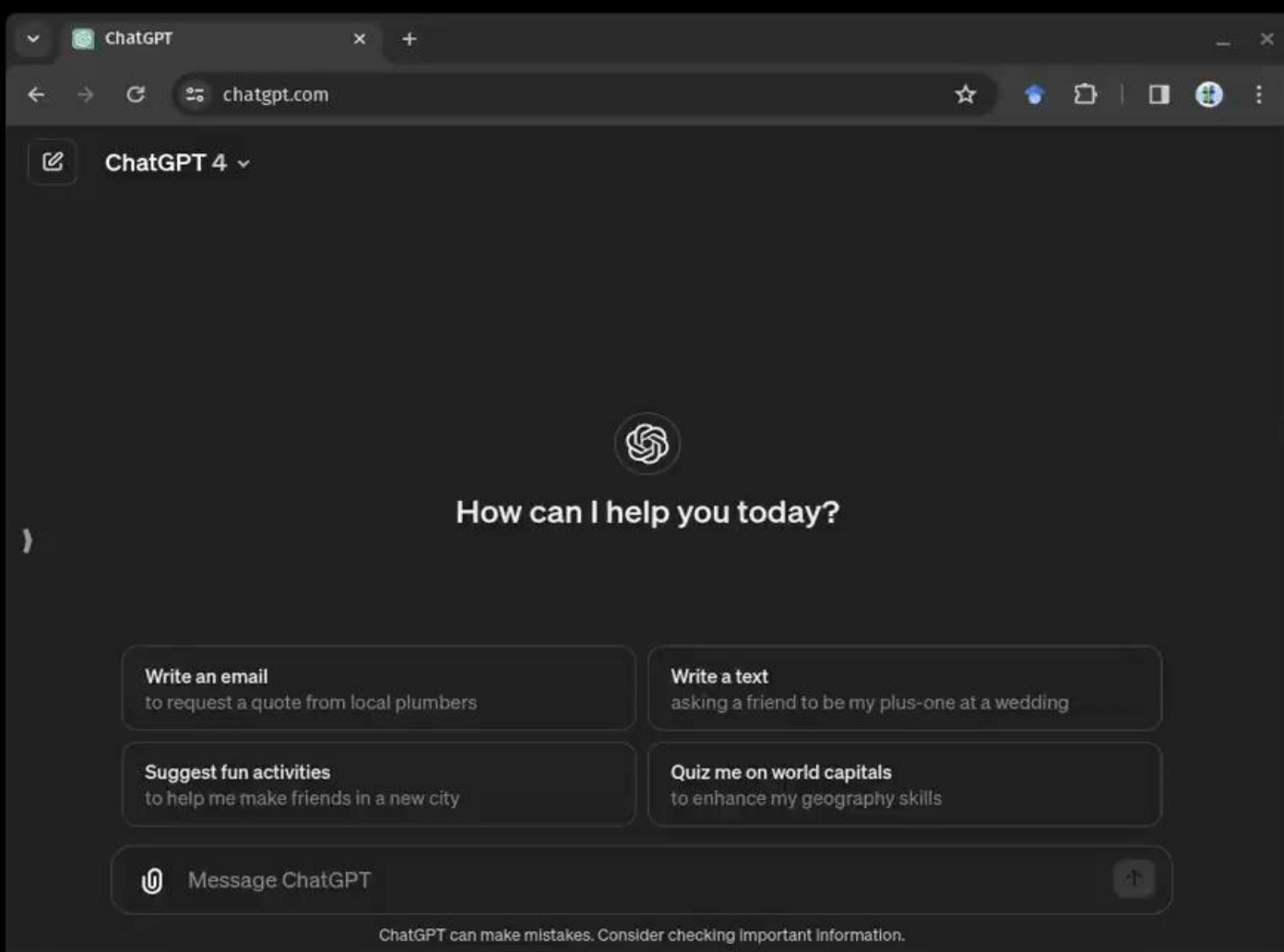




Specify

# From Code To Specs

Importance of  
legacy codes



**Generate**

# Decidable First-order Theories

Hoare logic,  
Separation logic

Automated Theorem  
Provers (ATP), e.g.  
Dafny, Frama-C, Verus

```
method maxArray(a: array<int>) returns (m: int)
  requires a != null && a.Length > 0
  ensures forall k :: 0 <= k < a.Length ==> a[k] <= m
  ensures exists k :: 0 <= k < a.Length && a[k] == m
{
  m := a[0];
  var i := 1;
  while (i < a.Length)
    invariant 1 <= i <= a.Length
    invariant forall j :: 0 <= j < i ==> a[j] <= m
    invariant exists j :: 0 <= j < i && a[j] == m
  {
    m := if m > a[i] then m else a[i];
    i := i + 1;
  }
}
```

# Expressive Higher-Order Theories

Dependent  
type theory

Interactive Theorem  
Provers (ITP), e.g.  
Coq, Lean, Agda

```
-- Import necessary libraries for proving
import data.nat.basic
import tactic

-- Using the previously defined factorial function
def factorial :  $\mathbb{N} \rightarrow \mathbb{N}$ 
| 0      := 1
| (n+1) := (n + 1) * factorial n

-- Proof that factorial is always positive
theorem factorial_pos :  $\forall n : \mathbb{N}, \text{factorial } n > 0 :=$ 
begin
  -- Apply mathematical induction on n
  apply nat.rec,
  -- Base case: show that factorial 0 is 1, which is positive
  { show factorial 0 > 0,
    rw factorial, -- factorial 0 = 1 by definition
    exact nat.zero_lt_one, },
  -- Inductive step: assume factorial n is positive,
  -- prove factorial (n+1) is positive
  { intros n h,
    show factorial (n + 1) > 0,
    rw factorial, -- expand factorial (n + 1)
    apply nat.mul_pos, -- product of two positive numbers is positive
    show n + 1 > 0,
    exact nat.succ_pos n,
    exact h, }
end
```

# Refinement-based Synthesis

*Functionality* without  
sacrificing *performance*

e.g. fiat-crypto library:

Estimated that over 95% of  
HTTPS connections by browsers  
run the generated algorithm.

$$\text{insert } k \ v \ I \equiv \{I' \mid I' \subseteq [(k, v)] \cup I\}$$

$$\downarrow \text{IU}$$

$$\{I' \mid k \notin I \rightarrow I' \subseteq [(k, v)] \cup I \\ \wedge k \in I \rightarrow I' \subseteq I\}$$

$$\downarrow \text{IU}$$

$$\{I' \mid k \notin I \rightarrow I' = [(k, v)] \cup I \\ \wedge k \in I \rightarrow I' \subseteq I\}$$

$$\downarrow \text{IU}$$

$$b \leftarrow \{b \mid \text{if } b \text{ then } k \notin I \text{ else } k \in I\};$$

$$\text{if } b \text{ then ret } [(k, v)] \cup I \text{ else } \{I' \mid I' \subseteq I\}$$

$$\downarrow \text{IU}$$

$$b \leftarrow \text{ret notKey}(k, I);$$

$$\text{if } b \text{ then ret } [(k, v)] \cup I \text{ else ret } I$$

$$\downarrow \text{IU}$$

$$\text{if notKey}(k, I) \text{ then ret } [(k, v)] \cup I \\ \text{else ret } I$$

# Correct By Construction

Tactic-driven

Code-carrying proof

Extracting implementations

## Insight 1

Construct tactic by tactic,  
not sample token by token

```
hone method insert. {  
  StartMethod.  
  setoid_rewrite refine_ReplaceUsedKeyAdd.  
  setoid_rewrite refine_SubEnsembleInsert.  
  autorewrite with monad laws.  
  setoid_rewrite refine_pick_KeyToBeReplaced min.  
  setoid_rewrite refine_If_Then_Else_Bind.  
  autorewrite with monad laws.  
  setoid_rewrite refine_If_Opt_Then_Else_Bind.  
  autorewrite with monad laws.  
  setoid_rewrite refine_pick_CacheADTwLogIndex AbsR.  
  setoid_rewrite refine_pick_KVEnsembleInsertRemove  
    with (1 := EquivKeys H).  
  setoid_rewrite refine_pick_KVEnsembleInsert  
    with (1 := EquivKeys H).  
  autorewrite with monad laws; simpl.  
  finish honing. }
```

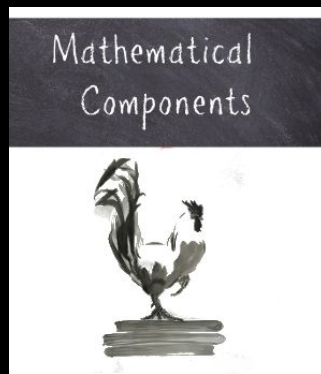
# Correct By Construction

## Insight 2

For AI models to generalize well, we need **generic tactics**

Develop tactics for categories – more precisely, algebraic theories in the form of **typeclasses**

e.g. proof of Four Color Thm and Odd Order Thm with mathematical components



```
Definition eq_axiom T (e : rel T) := forall x y, reflect (x = y) (e x y).
```

```
HB.mixin Record hasDecEq T := { eq_op : rel T; eqP : eq_axiom eq_op }.
```

```
#[mathcomp(axiom="eq_axiom"), short(type="eqType")]
```

```
HB.structure Definition Equality := { T of hasDecEq T }.
```

```
Lemma eq_refl (T : eqType) (x : T) : x == x. Proof. exact/eqP. Qed.
```

```
Notation eqxx := eq_refl.
```

```
Lemma eq_sym (T : eqType) (x y : T) : (x == y) = (y == x).
```

```
Proof. exact/eqP/eqP. Qed.
```

```
#[global] Hint Resolve eq_refl eq_sym : core.
```

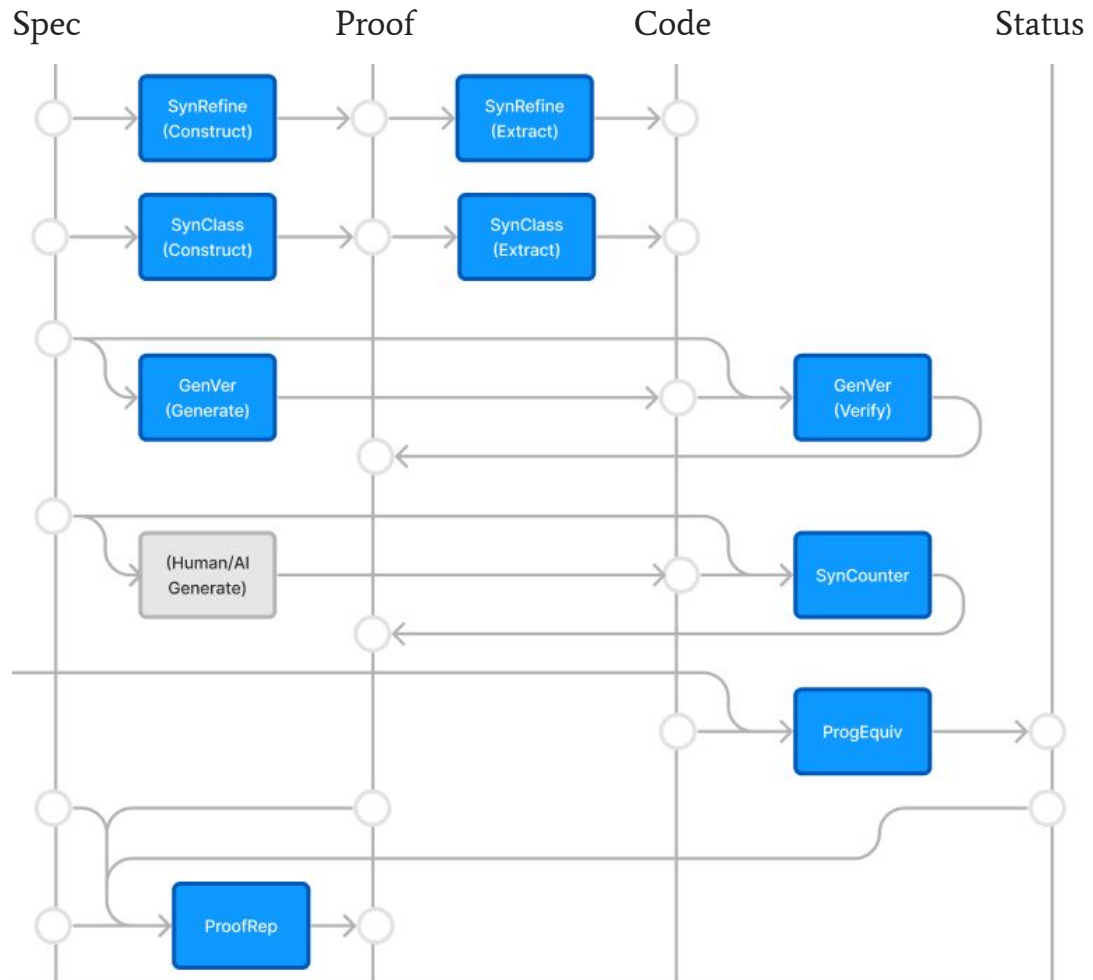
## Generate

### Program Equivalence

- Check if two implementations have the same behavior (ProgEquiv)
- Otherwise give a counterexample (ProgCounter)

### Proof Repair

- Fix proofs when libraries, specs, implementations change (ProofRep)





# Summary

<b>Challenge</b>	Scaling formal verification with AI assistance
<b>Model</b>	Feeding AI models with legacy code and doc
<b>Interact</b>	Human users + proof assistant + AI assistant
<b>Specify</b>	From natural language to formal structures
<b>Generate</b>	Correct by construction, not by generation

**Thank you!**

# Lean Copilot

suggest\_tactics

search\_proof

select\_premises

The screenshot shows the Lean IDE interface with a file named `Lean4Example.lean`. The code contains two theorems:

```
1 -- Lean Copilot is open-sourced at https://github.com/lean-dojo/LeanCopilot
2 import LeanCopilot
3 import Mathlib.Data.Set.Basic
4
5
6 theorem add_abc : ∀ a b c : N, a + b + c = a + c + b := by
7   intro a b c
8   simp [Nat.add_assoc, Nat.add_comm b]
9
10
11 theorem set_inter_comm (s t : Set α) : s ∩ t = t ∩ s := by
12   ext x
13   simp_all only [Set.mem_inter_iff]
14   apply Iff.intro
15   · intro a
16     simp_all only [and_self]
17   · intro a
18     simp_all only [and_self]
19
```

The right-hand pane shows the `Lean Infoview` with the following content:

- Lean4Example.lean:19:0
- Tactic state
- No goals
- All Messages (0)

The status bar at the bottom indicates the current position: `Ln 19, Col 1`, `Spaces: 2`, `UTF-8`, `LF`, and `lean4`.

# Dependencies

